

量子情報科学ウィンタースクール 量子秘密分散法

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What is Quantum Secret Sharing Schemes ?

○ Classical Secret Sharing Schemes (SSS)

- (k, n) -threshold SSS (Shamir 1979, Blakley 1979)
- (k, L, n) -threshold ramp SSS
(Yamamoto 1985, Blakley-Meadows 1985)

○ Quantum Secret Sharing Schemes (QSSS)

$\left\{ \begin{array}{l} \text{to encode classical messages (bit)} \\ \text{to encode quantum states (qbit)} \end{array} \right\}$ into quantum states

○ What for ?

- outputs of quantum computer
- outputs of expensive apparatus
- quantum key in cryptography with quantum algorithm

Literature on QSSS

- (k, n) -threshold QSSS (Cleve-Gottesman-Lo, 1999)
- Coding efficiency of perfect QSSS (Gottesman, 2000)
- Information theoretical treatment (Imai et al., 2003)
based on **coherent information** and reference system

Our Goal

- Information theoretical treatment
based on **reversibility** and **Holevo Information**
- Coding efficiency of **ramp QSSS** and optimal construction

Quantum Secret Sharing Schemes (QSSS)

○ $\mathcal{S}(\mathcal{H})$: totality of density operators on \mathcal{H}

\mathcal{H} : Hilbert space

ρ : original state $\in \mathcal{S}(\mathcal{H})$

\Downarrow W_N : encoder into shares $N = \{1, 2, \dots, n\}$

$\mathcal{H}_N = \mathcal{H}_1 \otimes \mathcal{H}_2 \otimes \dots \otimes \mathcal{H}_n$

$W_N(\rho)$: entangled state

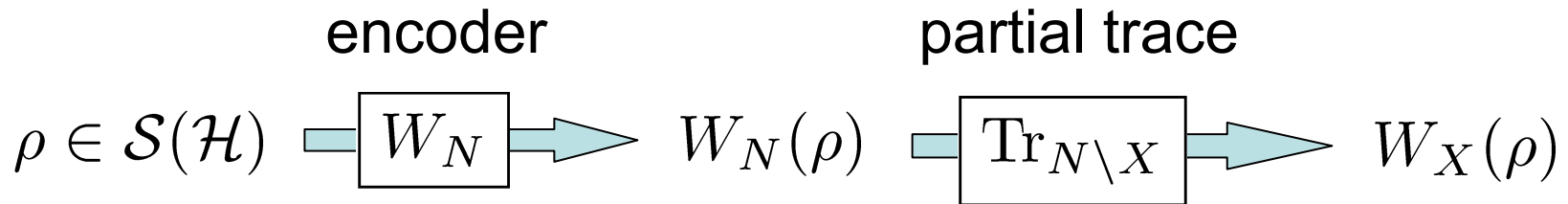
partial trace for a subset $X \subseteq N$

$\mathcal{H}_X = \bigotimes_{i \in X} \mathcal{H}_i$

$W_X(\rho) = \text{Tr}_{N \setminus X} \cdot W_N(\rho)$

$W_X(\rho)$ reproduces ρ or not ?

Authorized and Unauthorized Sets



○ subset $X \subseteq N$ $W_X = \text{Tr}_{N \setminus X} \cdot W_N$

○ X : authorized

$\stackrel{\text{def}}{\iff} W_X(\rho)$ can reproduce ρ for $\forall \rho \in \mathcal{S}(\mathcal{H})$

○ X : unauthorized

$\stackrel{\text{def}}{\iff} W_X(\rho)$ has no information about $\rho \in \mathcal{S}(\mathcal{H})$ for $\forall \rho$

○ X : intermediate $\stackrel{\text{def}}{\iff}$ otherwise

Perfect Schemes and Ramp Schemes

$$\rho \in \mathcal{S}(\mathcal{H}) \xrightarrow{W_N} W_N(\rho) \xrightarrow{\text{Tr}_{N \setminus X}} W_X(\rho)$$

- QSSS W_N : perfect scheme

$\stackrel{\text{def}}{\iff} X$ is either authorized or unauthorized for $\forall X \subseteq N$

- QSSS W_N : ramp scheme

$\stackrel{\text{def}}{\iff}$ otherwise

Reversibility of Quantum Operations

- quantum operation

$$\rho \in \mathcal{S}(\mathcal{H}) \xrightarrow{\boxed{\mathcal{E}}} \mathcal{E}(\rho)$$

{ trace preserving
completely positive
affine map

- subset $\mathcal{S} \subseteq \mathcal{S}(\mathcal{H})$

- $\mathcal{E} : \text{reversible w.r.t. } \mathcal{S}$

$\stackrel{\text{def}}{\iff} \exists \mathcal{R} \text{ such that}$

$$\forall \rho \in \mathcal{S} \xrightarrow{\boxed{\mathcal{E}}} \mathcal{E}(\rho) \xrightarrow{\boxed{\mathcal{R}}} \mathcal{R}\mathcal{E}(\rho) = \rho$$

- $\mathcal{E} : \text{vanishing w.r.t. } \mathcal{S}$

$\stackrel{\text{def}}{\iff} \exists \rho_0 \text{ such that}$

$$\forall \rho \in \mathcal{S} \xrightarrow{\boxed{\mathcal{E}}} \mathcal{E}(\rho) = \rho_0$$

Quantum Relative Entropy

○ $D(\rho||\sigma) \stackrel{\text{def}}{=} \text{Tr}[\rho(\log \rho - \log \sigma)]$

○ **monotonicity**

$$\begin{array}{ccc}
 \rho & \xrightarrow{\quad \mathcal{E} \quad} & \mathcal{E}(\rho) \\
 \sigma & & \mathcal{E}(\sigma)
 \end{array}$$

$$D(\rho||\sigma) \geq D(\mathcal{E}(\rho)||\mathcal{E}(\sigma))$$

— Theorem (Petz, 1986) —

The following conditions are equivalent.

1. \mathcal{E} is reversible w.r.t. $\{\rho, \sigma\}$
2. $D(\rho||\sigma) = D(\mathcal{E}(\rho)||\mathcal{E}(\sigma))$

Holevo Information

- \mathcal{E} : quantum operation
- μ : probability measure on $\mathcal{S} \subseteq \mathcal{S}(\mathcal{H})$

$$\mathbb{E}_\mu[\cdot] \stackrel{\text{def}}{=} \int_{\mathcal{S}} \cdot \mu(d\rho) \quad \sigma_\mu \stackrel{\text{def}}{=} \mathbb{E}_\mu[\rho]$$

- Holevo Information for \mathcal{E} and μ

$$I(\mu; \mathcal{E}) \stackrel{\text{def}}{=} \mathbb{E}_\mu[D(\mathcal{E}(\rho) || \mathcal{E}(\sigma_\mu))]$$

- also written as

$$I(\mu; \mathcal{E}) = H(\mathcal{E}(\sigma_\mu)) - \mathbb{E}_\mu[H(\mathcal{E}(\rho))]$$

$$H(\rho) \stackrel{\text{def}}{=} -\text{Tr}[\rho \log \rho] : \text{ von Neumann entropy}$$

Holevo Information and Reversibility

- $\mathcal{P}_+(\mathcal{S})$: set of probability measures on $\mathcal{S} \subseteq \mathcal{S}(\mathcal{H})$
- $\mu \in \mathcal{P}_+(\mathcal{S})$
- **monotonicity** $I(\mu; \mathcal{I}) \geq I(\mu; \mathcal{E})$ (\mathcal{I} : identity)

Theorem 1

The following conditions are equivalent.

1. \mathcal{E} is **reversible** (resp. **vanishing**) w.r.t. \mathcal{S}
2. $\forall \mu \in \mathcal{P}_+(\mathcal{S}), I(\mu; \mathcal{E}) = I(\mu; \mathcal{I})$ (resp. $= 0$)
3. $\exists \mu \in \mathcal{P}_+(\mathcal{S}), I(\mu; \mathcal{E}) = I(\mu; \mathcal{I})$ (resp. $= 0$)

- cf. Schumacher and Nielsen, 1996

Authorized (Unauthorized) Condition for Shares

- $\mathcal{E} = W_X$ $\mathcal{S} = \mathcal{S}_1(\mathcal{H})$: set of pure states
- reversibility w.r.t. $\mathcal{S}(\mathcal{H}) \iff$ reversibility w.r.t. $\mathcal{S}_1(\mathcal{H})$
- Holevo information for $\mu \in \mathcal{P}_+(\mathcal{S}_1(\mathcal{H}))$

$$I(\mu; \mathcal{I}) = H(\sigma_\mu) - \underbrace{\mathbb{E}_\mu[H(\rho)]}_{\text{0 for pure state}} \geq I(\mu; W_X)$$

0 for pure state

Theorem 2

The following conditions are equivalent.

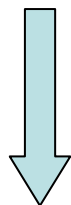
1. X is authorized (resp. unauthorized)
2. $\forall \mu \in \mathcal{P}_+(\mathcal{S}_1(\mathcal{H})), I(\mu; W_X) = H(\sigma_\mu)$ (resp. $= 0$)
3. $\exists \mu \in \mathcal{P}_+(\mathcal{S}_1(\mathcal{H})), I(\mu; W_X) = H(\sigma_\mu)$ (resp. $= 0$)

Coding Efficiency of Perfect QSSS

Theorem 3 (cf. Imai et al. 2003)

$$\forall X \subseteq N \quad \forall \mu \in \mathcal{P}_+(\mathcal{S}_1(\mathcal{H}))$$

$$H(\sigma_\mu) \leq H(W_X(\sigma_\mu))$$



μ : invariant measure on $\mathcal{S}_1(\mathcal{H})$

(uniform distribution on pure states)

$$\implies \sigma_\mu = I / \dim \mathcal{H}$$

Corollary 1 (Gottesman 2000)

$$\forall i \in N \quad \dim \mathcal{H} \leq \dim \mathcal{H}_i$$

Proof of Theorem 3 (1)

$$\begin{aligned} \bigcirc \text{ For } X \subseteq N \quad \exists Y : \text{unauthorized} &\implies I(\mu; W_Y) = 0 \\ X \cup Y : \text{authorized} &\implies I(\mu; W_{XY}) = H(\sigma_\mu) \end{aligned}$$

$$\begin{aligned} \bigcirc H(\sigma_\mu) &= I(\mu; W_{XY}) - I(\mu; W_Y) \\ &= H(W_{XY}(\sigma_\mu)) - \mathbf{E}_\mu[H(W_{XY}(\rho))] - H(W_Y(\sigma_\mu)) + \mathbf{E}_\mu[H(W_Y(\rho))] \\ &\leq H(W_X(\sigma_\mu)) - \mathbf{E}_\mu[H_\rho(W_X|W_Y)] \end{aligned}$$

\bigcirc subadditivity

$$H(W_{XY}(\sigma_\mu)) \leq H(W_X(\sigma_\mu)) + H(W_Y(\sigma_\mu))$$

\bigcirc conditional entropy

$$H_\rho(W_X|W_Y) \stackrel{\text{def}}{=} H(W_{XY}(\rho)) - H(W_Y(\rho))$$

Proof of Theorem 3 (2)

○ Y : unauthorized $X \cup Y$: authorized

$$H(\sigma_\mu) \leq H(W_X(\sigma_\mu)) - \mathbb{E}_\mu[H_\rho(W_X|W_Y)]$$

○ classical case : $H_\rho(W_X|W_Y) \geq 0$

○ quantum case : ~~$H_\rho(W_X|W_Y) \geq 0$~~

Proof of Theorem 3 (3)

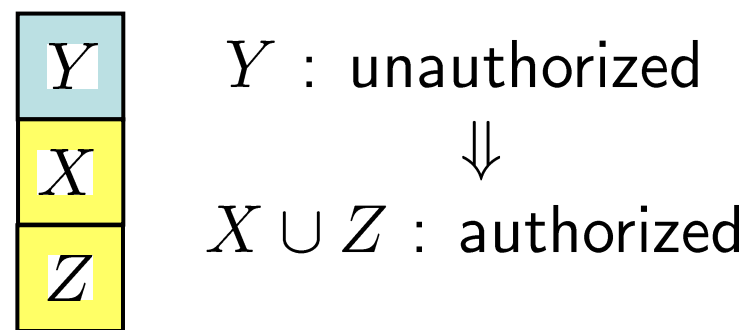
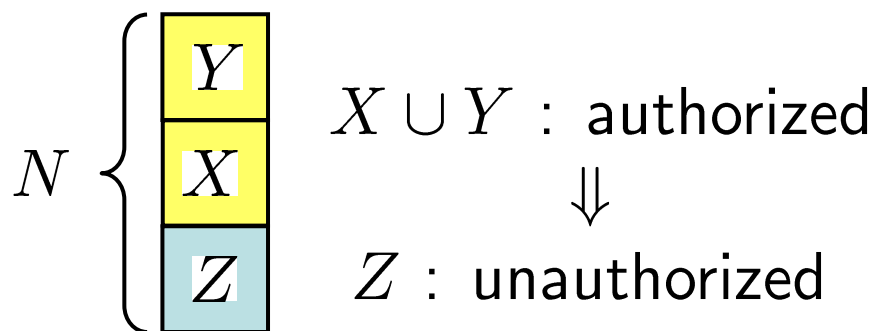
- ρ : pure state $\implies W_N(\rho)$: pure state
(\because Steinspring dilation, called pure state scheme)

- $X \cap Y = \emptyset$

$$Z \stackrel{\text{def}}{=} N \setminus X \cup Y$$

- no cloning theorem

- no deleting theorem



- Z has the same property as Y !

Proof of Theorem 3 (4)

$$\begin{cases} Y : H(\sigma_\mu) \leq H(W_X(\sigma_\mu)) - \mathbb{E}_\mu[H_\rho(W_X|W_Y)] \\ Z : H(\sigma_\mu) \leq H(W_X(\sigma_\mu)) - \mathbb{E}_\mu[H_\rho(W_X|W_Z)] \end{cases}$$

$$\Downarrow$$

$$H(\sigma_\mu) \leq H(W_X(\sigma_\mu)) - \frac{1}{2} \{ \mathbb{E}_\mu[H_\rho(W_X|W_Y)] + \mathbb{E}_\mu[H_\rho(W_X|W_Z)] \}$$

$$= H(W_X(\sigma_\mu))$$

0 for pure state scheme